

osmo-gmr: What's up with sat-phones ?

Piecing together the missing bits

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Introduction

Outline

- 1 GMR Introduction
- 2 GMR-1 Speech codec
- 3 GMR-1 Cipher

About the speaker

- Linux and free software "geek" since 1999
- M.Sc. in C.S. + some E.E.
- General orientation towards low level
 - Embedded, Kernel, Drivers and such.
 - Hardware (Digital stuff, FPGA or RF)
- Interest in various telecom and SDR projects for several years
 - Osmocom projects (OpenBSC, Osmocom-BB, ...)
 - Airprobe, OpenBTS ...
 - In my spare time

GMR Introduction

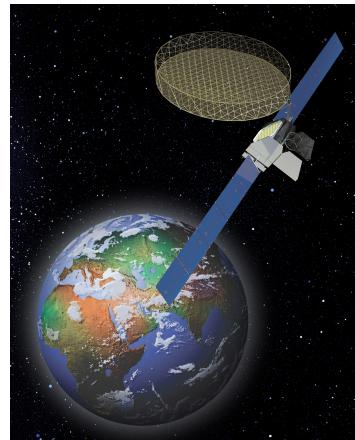
What is GMR ?

- "GEO-Mobile Radio Interface"
(GEO stands for Geostationary Earth Orbit)
- ETSI standard for satellite phones
- Heavily based on GSM
- Multiple standards :
 - GMR-1 (ETSI TS 101 376)
 - GMR-1 (*the one described in this talk*)
 - GmPRS
 - GMR-1 3G
 - GMR-2 (ETSI TS 101 377)

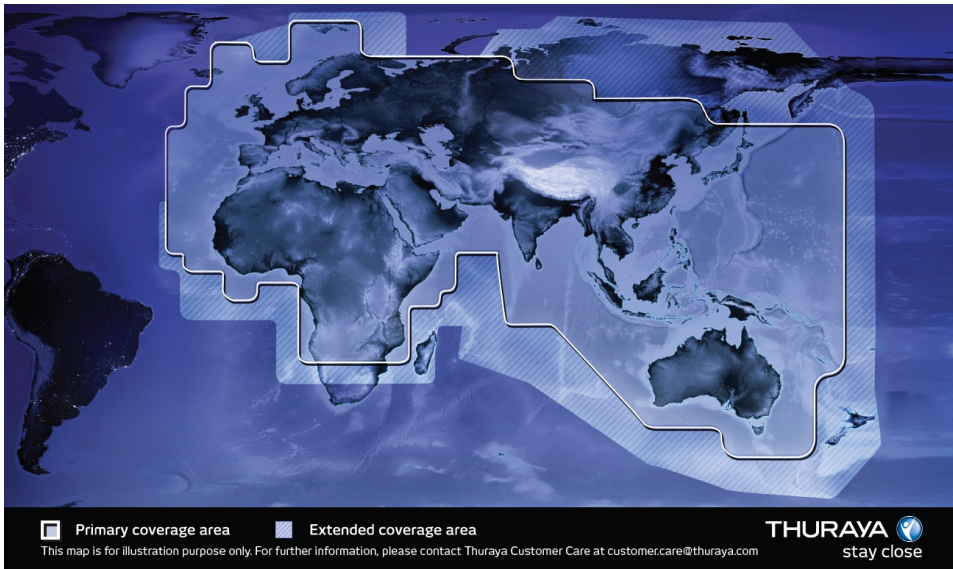


Deployment

- GMR-1
 - Thuraya
 - Thuraya 2 (44E) and Thuraya 3 (98.5E)
 - Main focus of our attention so far
 - SkyTerra ?
 - TerreStar ?
 - ICO
 - (Inmarsat R-BGAN)
 - Solaris Mobile (*future*)
- GMR-2
 - Inmarsat "IsatPhone"
 - ACes



Deployment Thuraya



Comparison to GSM

Features

- New names
 - BTS → GTS, BSC → GSC, BSS → GSS, ...
 - MS → MES(-MS)
- New Specialized features
 - Terminal-to-Terminal calls
 - High Penetration Alerting (HPA)
- Tight links to GPS
 - Almanac and Ephemeris sent by the satellite
 - Position reported in RACH (Channel Request)
- New speech codec: AMBE
- New cipher

Comparison to GSM

Protocol Stack

- Layer 0/1: Completely different
 - Different bursts and TDMA multiplex / multi-frame
 - Different modulation
 - More channels types
- Layer 2: LAPSat vs LAPDm
 - Both simplified version of LAPD
 - Shorter header
 - k=16 window size for outstanding unacknowledged segments
- Layer 3:
 - RR different
 - MM/CM common
- Same core network
- Packet Data:
 - RLC/MAC different
 - LLC and above common/shared

Previous work

osmo-gmr

- Hardware setup
 - Antennas, LNAs, filters, SDR receivers
- SDR processing
 - Channelization, Demodulation
- Channel coding
 - Channel types, Interleaving, Viterbi, CRC
- Demo receive app
 - Synchronization, minimal higher level support
- Wireshark dissector
- Two big missing pieces for basic RF to Audio chain
 - Voice Codec
 - Cipher

Refer to the 28C3 talk for more details

GMR-1 Speech codec

GMR-1 Speech codec

The problem

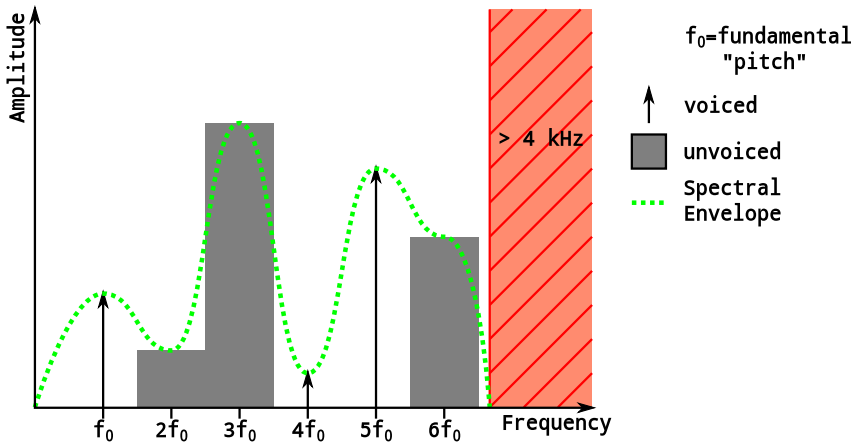
- AMBE: Advanced Multi-Band Excitation
- Not documented in the standard
 - Barely a high level description
 - No reference code
- Proprietary codec by *DVSI Inc.*
 - Not supported by their "cheap" hardware USB decoder
 - Cheapest hardware is the NET-2000 appliance (2kEUR)
- But :
 - mbelib: Code for other documented IMBE/AMBE variants (P25)
 - Implemented in SO2510 phone DSP (TI C55x)

AMBE Codec

Description

- Highly specialized for voice (vocoder)
- Divides speech in small segments
 - For GMR-1: 20 ms frames subdivided into co-quantized 10 ms sub-frames
- Represent each speech (sub)frame as a set of parameters
 - f_0 : Fundamental frequency (*pitch*)
 - G : Gain (*volume*)
 - Voiced / Unvoiced decision (per band)
 - Spectral Magnitudes
- Decoding can be summarized as 3 steps:
 - **Unpacking**: Unpack the raw frame bits into quantized parameters
 - **De-Quantization**: From quantized parameters to actual values
 - **Synthesis**: From the parameters set to actual audio

AMBE Codec Synthesis



AMBE Reversing

DSP Code analysis

- Target: SO-2510 phone
- Codec has to be in the DSP, nowhere else it could be !
- DSP firmware extracted from firmware update package
 - Supported by IDA
- But where ?
 - 250k binary blob
 - No strings
 - Obscure TI C55x assembly
- Dieter Spaar to the rescue !
 - Identified entry points for encode/decode functions
 - Look for Audio DMA / Interrupts
 - Search for constants
 - Stack Switching



AMBE Reversing Simulator

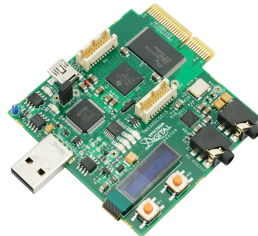
- TI Code Composer Studio - Simulator
 - Accurately simulates supported DSP
 - Arbitrary memory layout
 - fread()/fwrite() from host
 - Tracing of all memory access
 - Windows only :(
- Use the original firmware to decompress audio for us
 - DSP dump converted to a valid COFF .OBJ file for linking
 - Custom linker script
 - Simple main() that fread() frames and fwrite() audio
- Success !
 - It took quite a few tries, lots of traps
 - But it works and we get audio out
 - Slow (not real-time) and not practical though



AMBE Reversing

Hardware

- Real HW would be faster and more convenient. But :
 - Code has to run at the physical address it has been linked for
 - OMAP has a DSP MMU, but standalone DSP don't
 - Need a cheap board with a compatible memory map
- Dieter found one with SDRAM where needed and Ethernet
 - Success ! About 16x faster than real-time
 - SDRAM is not fast, relocate some data tables to SRAM
- I indented to buy the same board
 - But in my haste ... I ordered the wrong one ... ***facepalm***
 - No SDRAM, more SRAM, but at the wrong physical address
 - Easy, just relocate the code ! Can't be that hard, right ?
 - Use IDAPython + simulator trace mode
 - Success !



AMBE Reversing

Software

- Hardware USB decoder is nice, but not enough
- Decompression process:
 - **Unpacking**
 - Early, simple bit manipulation, easy to follow
 - **Dequantization**
 - Easily 95% of the work
 - Hard to follow fixed point math in DSP assembly
 - **Synthesis**
 - Started by just re-using mbelib code
 - Then rewrote using P25 specs and some guessing
- Resulting PoC/reference code in GIT
 - Not same audio quality as the original, but perfectly intelligible

```
amar    **+AR0(#857h)
call    sub_103540
mov     *AR6(#100h), AR1
btst    @0, AR1, TC1
mov     #0, T2
mov     xccpart !TC1
bcc     loc_10B148, !TC1
amar    *AR7, XAR0
amar    *AR5, XAR1
|| mov  T2, T0
call    sub_109EC8
mov     T0, T2
```

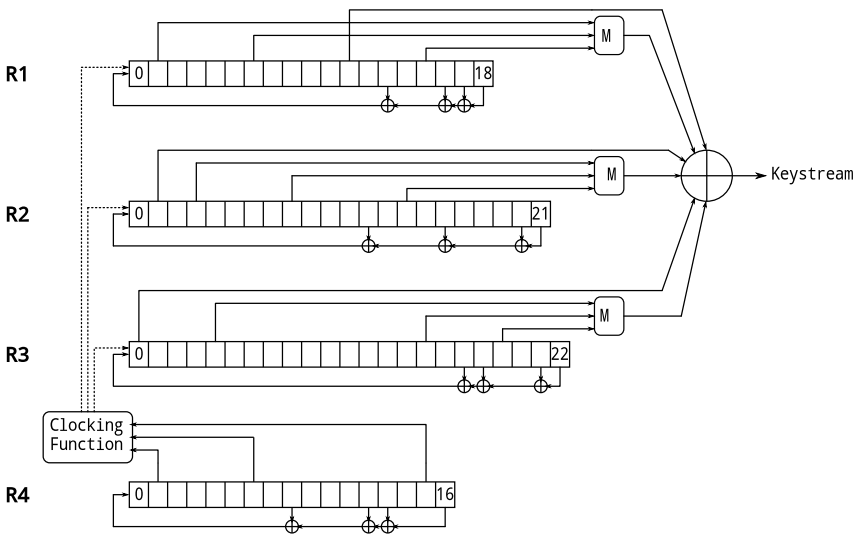
```
amar    *AR5, XAR3
amar    *AR5, XAR4
amar    **+AR3(#684h)
amar    *AR3, XAR2
amar    **+AR4(#684h)
amar    *AR5, XAR3
amar    **+AR3(#584h)
rpt     #0FFh
mov     *AR4+, *AR3+
mov     *AR5(#785h), AR1
mov     AR1, *AR5(#784h)
amar    *SP(#var_0), XAR3
|| rpt #0FFh
mov     *AR3+, *AR2+
add     #102h, mmap(@SP)
mov     T3, *AR5(#785h)
popboth XAR6
```

GMR-1 Cipher

Extraction from DSP

- Performed by a team at Bochum university (RUB)
 - Led by Benedikt Driessen
- Published beginning 2012
- Validated against on-the-air data using `osmo-gmr`
- Process:
 - Extract DSP image by running the ARM DSP download code in QEMU
 - Analyze DSP image using heuristics
 - (But really, if you look for XOR for a few minutes, that works too)
 - Refer to <http://gmr.crypto.rub.de> for more details and the papers
- Also includes a cipher-text only attack

A5-GMR-1 Structure



A5-GMR-1

Description

- Based off GSM's A5/2
 - Feedback, output and clocking taps changed
 - Output function and Initialization tweaked
 - But globally the same structure
- 4 maximal length LFSRs (R1=19, R2=22, R3=23 and R4=17)
- Initialization:
 - Mix the key and frame-number with a linear function
 - Clock all LFSRs 64 times while mixing-in the result from above
 - Force the MSB of each LFSRs to 1
 - Clock all LFSRs 250 times
- Bitstream generation:
 - Clock the cipher as many times as needed to generate enough bits
 - R4 is tapped into a clocking function that drives R1,R2,R3 clocking
 - R1, R2 and R3 are tapped and combined into an output bit

Linear algebra over $GF(2)$ quick refresher

- $GF(2)$ is a fancy term for *binary*
 - Addition is the logic XOR
 - Multiplication is the logic AND
- Linear: each term is a constant or a variable/unknown multiplied by a constant
- Linear equations systems :
 - Can be represented as matrix operations
$$\begin{array}{l} a \cdot x + b \cdot y = c \\ d \cdot x + e \cdot y = f \end{array} \iff \begin{pmatrix} a & b \\ d & e \end{pmatrix} \cdot \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} c \\ f \end{pmatrix}$$
 - Can be solved efficiently
 - Number of independent equations vs number of unknowns determines the number of possible solutions (over/under determination)
- Other linear operations can also be rewritten as matrix operations
 - State change of a LFSR for instance: $S_{t+1} = A \cdot S_t$
 - They can also be combined: $S_{t+n} = A^n \cdot S_t$
 - If they add redundancy (like FEC), a parity-check matrix can check the result

RUB attack

- Cipher-text only
- Targets TCH3 traffic/voice frame
- Based off previous A5/2 work
 - "Instant Ciphertext-Only Cryptanalysis of GSM encrypted communication" by E.Barkan, E.Biham, and N.Keller
 - Tweaked for GMR-1 and its TCH3 frames
 - Adds more "guessed" bits in each register to reduce the unknowns
- Results :
 - 350 Gb of off-line data
 - 32 TCH3 frames
 - About 40 minutes on-line phase

A better attack

Overview

- Based on the same A5/2 GSM attack
 - Don't do anything fancy, just tweak for A5-GMR-1
- Both known-plaintext and ciphertext-only variant
- Targets FACCH3 control frames instead of TCH3 voice frames
- FACCH3 advantages :
 - Simpler modulation and better training sequence → less bit-errors
 - Predictable plaintext → known-plaintext attacks
 - Much more redundancy (more FEC) → less bursts needed for ciphertext-only attacks
 - Used to negotiate TCH6/TCH9 channels → attack works for CSD/Fax

A better attack

Known plaintext

File Edit View Go Capture Analyze Statistics Telephony Tools Internals Help

Filter: `gsmtap.gmr1_chan_type == 18` Expression... Clear Apply Save

No.	Time	Source	Destination	Src Port	Dst Port	Protocol	Fn	Info
178	0.114724	127.0.0.1	127.0.0.1	60654	4729	LAPSat	22996 I	N(R)=4, N(S)=5 (DTAP) (RR) Ciphering Mode Command
183	0.119807	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23020 S	func=RR, N(R)=5
185	0.121513	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23024 S	func=RR, N(R)=6
187	0.123068	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23028 S	func=RR, N(R)=7
188	0.124150	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23032 I	N(R)=7, N(S)=6 (Fragment)
190	0.125702	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23036 I	N(R)=7, N(S)=7 (DTAP) (MM) MM Information
191	0.126760	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23040 I	N(R)=7, N(S)=8 (DTAP) (MM) Location Updating Accept
193	0.128321	127.0.0.1	127.0.0.1	60654	4729	LAPSat	23044 I	N(R)=7, N(S)=9 (DTAP) (RR) Channel Release

▶ Frame 183: 68 bytes on wire (544 bits), 68 bytes captured (544 bits)
 ▶ Ethernet II, Src: 00:00:00 00:00:00 (00:00:00:00:00:00), Dst: 00:00:00 00:00:00 (00:00:00:00:00:00)
 ▶ Internet Protocol Version 4, Src: 127.0.0.1 (127.0.0.1), Dst: 127.0.0.1 (127.0.0.1)
 ▶ User Datagram Protocol, Src Port: 60654 (60654), Dst Port: gsmtap (4729)
 ▶ GSM TAP Header, ARFCN: 0 (Downlink), TS: 13, Channel: FACCH3 (0)
 ▼ Link Access Procedure, Satellite channel (LAPSat)
 ▶ Address Field: 0x00
 ▶ Control Field: S, func=RR, N(R)=5 (0x281)
 ... 0000 = Payload last nibble: 0x00
 Length Field: 0

```

0000 00 00 00 00 00 00 00 00 00 00 00 08 00 45 00 .....E.
0010 00 36 ea c3 40 00 40 11 51 f1 7f 00 00 01 7f 00 .....6..@.@.Q.....
0020 00 01 ec ee 12 79 00 22 fe 35 02 04 0a 0d 00 00 .....y."...5.....
0030 00 00 00 00 59 ec 12 00 00 00 80 28 10 00 00 00 .....Y...{.....
0040 00 00 00 00 .....
  
```

File: "/tmp/tnt-locupd-decipherer..." Packets: 261 · Displayed: 33 · Marked: 3 · Load time: 0:00... Profile: Default

A better attack

Plaintext attack

- Goal is to describe cipher as a linear operation: $A \cdot x = b$
 - A = matrix describing cipher, x = internal state and b = cipher stream
 - Each row of A and b is a bit of the output
- Internal cipher state dependency on FN and Kc is linear
 - Possible to combine equations from different bursts at different FN
 - Can recover Kc from the state
- Non-linear elements:
 - Majority function: $\mathcal{M}(a, b, c) = a + b.c$
 - Introduces quadratic terms
 - Linearize by adding one new unknown for every possible quadratic term
 - 594 new unknowns
 - Irregular clocking depending on R4 value
 - R4 is 17 bits but one is forced to '1' at init. Small enough for brute force !
 - Assume a given value for R4
 - Repeat 65536 times

A better attack

R4 quick scan

- In $A_n \cdot x = b$, some equations are redundant
- We can get a parity-check matrix H_n such that $H_n \cdot b = \mathbf{0}$
- Those 65536 H_n matrices can be precomputed offline
- With a single matrix-multiply we can check if a given R4 value is even a possibility
 - If result is non-zero, we can skip that R4 value
 - If result is zero, then we try to solve the system
 - In practice, only a few R4 value ever matches

A better attack

Ciphertext only

- Channel coding operation: $m = d \cdot G + g$
- Let H be the parity-check matrix so that $H \cdot (m + g) = \mathbf{0}$

- Encryption operation: $y = m + b$

- H can be used to derive equations from the ciphertext y :

$$\begin{aligned} H \cdot (y + g) &= H \cdot (m + b + g) \\ &= H \cdot b + \underbrace{H \cdot (m + g)}_0 \\ &= H \cdot A \cdot x \end{aligned}$$

- The same R4 quick-scan technique can also be used here
- To get enough equations for a unique solutions, multiple frames are needed

A better attack

Results

- Known-plaintext variant
 - Requires between 4 and 8 bursts depending on alignment
 - Space: 50 Mb
 - Time: 500 ms
- Ciphertext-only variant
 - Requires 8 consecutive bursts belonging to 2 FACCH3 L2 frames
 - Space: 5 Gb
 - Time: 1 s

Final words

Future

- C-band
- Packet Data (GmPRS)
- Upper layers implementation
- CSN.1 and 04.008 code generators
- TX side

Help welcome :)

Other satellite phone systems

- We choose Thuraya because :
 - Visible from Europe
 - Cheapest sat phone on ebay
 - Specifications mostly available
- Don't think other are better without proof
 - Availability of commerical intercepts tend to say otherwise

Thanks

Thanks to anyone who contributed to this projects and related ones. Most notably:

- Dimitri "horizon" Stolnikov
- Dieter Spaar
- RUB team

Further reading

■ GMR-1 in general

OsmocomGMR <http://gmr.osmocom.org/>

28C3 talk <http://gmr.osmocom.org/trac/blog/28c3-recording>

GMR1 Specs <http://www.etsi.org/standards-search>

GSM Specs <http://webapp.etsi.org/key/queryform.asp>

■ AMBE Codec

DVSI Inc. <http://www.dvsinc.com/>

■ GMR-1 Cipher

RUB GMR page <http://gmr.crypto.rub.de/>

Paper <http://cryptome.org/gsm-crack-bbk.pdf>